CSC D70:
Compiler Optimization
Parallelization

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The content of this lecture is adapted from the lectures of Todd Mowry and Tarek Abdelrahman

Announcements

- Final exam: Monday, **April 15**, **2:00-4:00pm**; Room: **IC200**
- Covers the whole semester
- Course evaluation (right now)

Data Dependence

$$
\begin{aligned}
S_1 &: \quad A = 1.0 \\
S_2 &: \quad B = A + 2.0 \\
S_3 &: \quad A = C - D \\
& \qquad \vdots \\
S_4 &: \quad A = B/C
\end{aligned}
$$

We define four types of data dependence.

- Flow (true) dependence: a statement S_i precedes a statement S_j in execution and S_i computes a data value that S_j uses.
- Implies that S_i must execute before S_j.

$$S_i \, \delta^t \, S_j \qquad (S_1 \, \delta^t \, S_2 \quad \text{and} \quad S_2 \, \delta^t \, S_4)$$

Data Dependence

$$
\begin{array}{ll}
S_1: & A = 1.0 \\
S_2: & B = A + 2.0 \\
S_3: & A = C - D \\
 & \quad \vdots \\
S_4: & A = B/C
\end{array}
$$

We define four types of data dependence.

- Anti dependence: a statement S_i precedes a statement S_j in execution and S_i uses a data value that S_j computes.
- It implies that S_i must be executed before S_j.

$$S_i \, \delta^a \, S_j \qquad (S_2 \, \delta^a \, S_3)$$

Data Dependence

$$
\begin{aligned}
S_1 &: \quad A = 1.0 \\
S_2 &: \quad B = A + 2.0 \\
S_3 &: \quad A = C - D \\
& \qquad \vdots \\
S_4 &: \quad A = B / C
\end{aligned}
$$

We define four types of data dependence.

- Output dependence: a statement S_i precedes a statement S_j in execution and S_i computes a data value that S_j also computes.
- It implies that S_i must be executed before S_j.

$$S_i \, \delta^o \, S_j \qquad (S_1 \, \delta^o \, S_3 \quad \text{and} \quad S_3 \, \delta^o \, S_4)$$

Data Dependence

$$
\begin{array}{ll}
S_1: & A = 1.0 \\
S_2: & B = A + 2.0 \\
S_3: & A = C - D \\
 & \quad\vdots \\
S_4: & A = B/C
\end{array}
$$

We define four types of data dependence.

- Input dependence: a statement S_i precedes a statement S_j in execution and S_i uses a data value that S_j also uses.
- Does this imply that S_i must execute before S_j?

$$S_i \, \delta^I \, S_j \qquad (S_3 \, \delta^I \, S_4)$$

Data Dependence (continued)

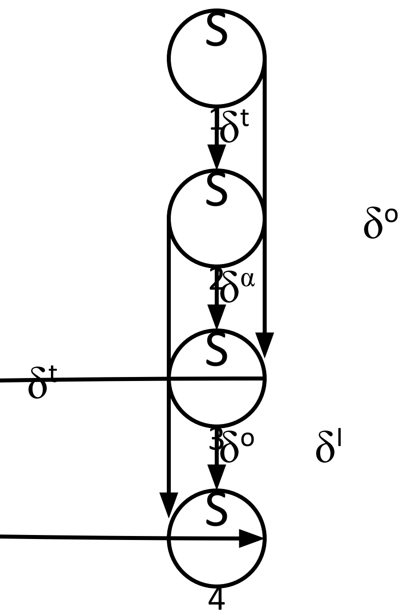
- The dependence is said to flow from S_i to S_j because S_i precedes S_j in execution.
- S_i is said to be the source of the dependence. S_j is said to be the sink of the dependence.
- The only "true" dependence is flow dependence; it represents the flow of data in the program.
- The other types of dependence are caused by programming style; they may be eliminated by re-naming.

$$
\begin{array}{ll}
S_1: & A = 1.0 \\
S_2: & B = A + 2.0 \\
S_3: & A1 = C - D \\
 & \vdots \\
S_4: & A2 = B/C
\end{array}
$$

Data Dependence (continued)

- Data dependence in a program may be represented using a dependence graph G=(V,E), where the nodes V represent statements in the program and the directed edges E represent dependence relations.

$$
\begin{aligned}
S_1 &: \quad A = 1.0 \\
S_2 &: \quad B = A + 2.0 \\
S_3 &: \quad A = C - D \\
&\qquad \vdots \\
S_4 &: \quad A = B/C
\end{aligned}
$$



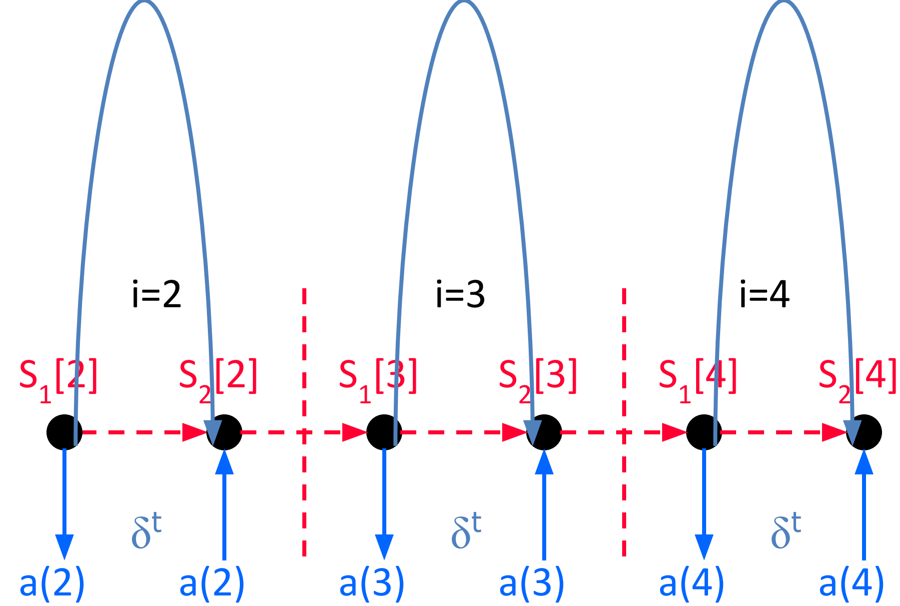
Value or Location?

- There are two ways a dependence is defined: value-oriented or location-oriented.

$$
\begin{aligned}
S_1 &: \quad A = 1.0 \\
S_2 &: \quad B = A + 2.0 \\
S_3 &: \quad A = C - D \\
&\quad\quad \vdots \\
S_4 &: \quad A = B/C
\end{aligned}
$$

Example 1

```
      do i = 2, 4
S1:     a(i) = b(i) + c(i)
S2:     d(i) = a(i)
      end do
```

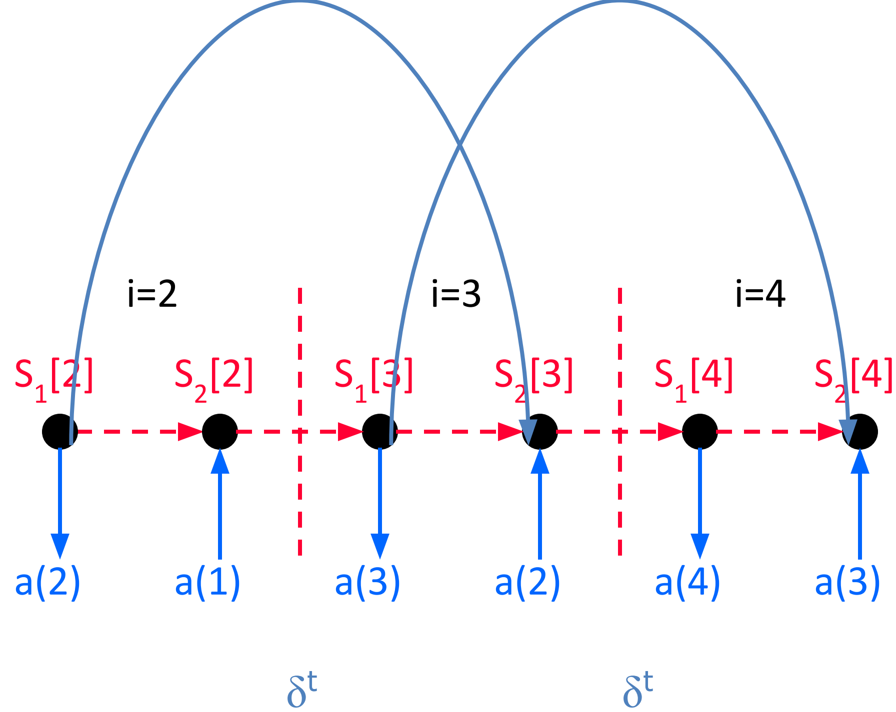


- There is an instance of S_1 that precedes an instance of S_2 in execution and S_1 produces data that S_2 consumes.
- S_1 is the source of the dependence; S_2 is the sink of the dependence.
- The dependence flows between instances of statements in the same iteration (loop-independent dependence).
- The number of iterations between source and sink (dependence distance) is 0. The dependence direction is =.

$$S_1\,\delta^t_=\,S_2 \qquad \text{or} \qquad S_1\,\delta^t_0\,S_2$$

Example 2

```
    do i = 2, 4
S1:   a(i) = b(i) + c(i)
S2:   d(i) = a(i-1)
    end do
```

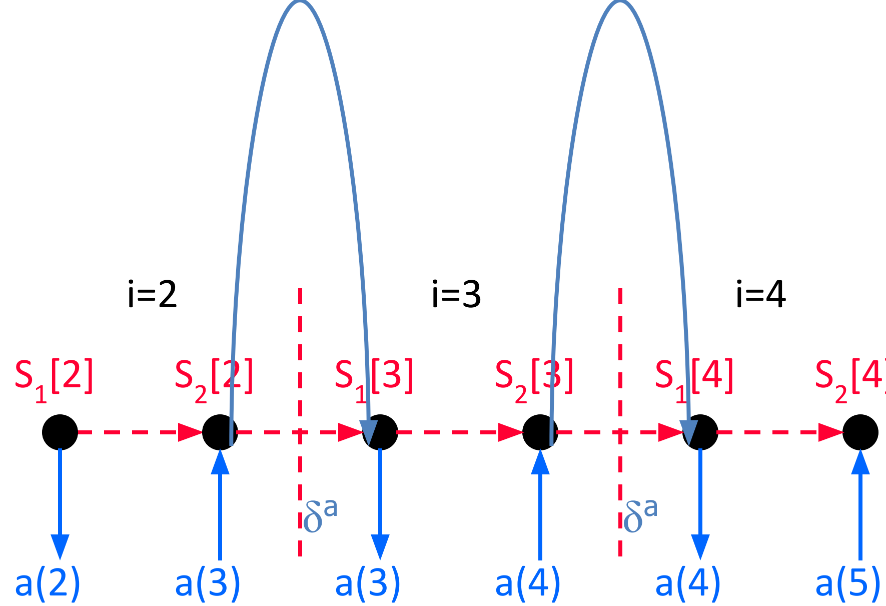


- There is an instance of S_1 that precedes an instance of S_2 in execution and S_1 produces data that S_2 consumes.
- S_1 is the source of the dependence; S_2 is the sink of the dependence.
- The dependence flows between instances of statements in different iterations (loop-carried dependence).
- The dependence distance is 1. The direction is positive (<).

$$S_1\, \delta^t_< \, S_2 \qquad \text{or} \qquad S_1\, \delta^t_1\, S_2$$

Example 3

do i = 2, 4
S_1: a(i) = b(i) + c(i)
S_2: d(i) = a(i+1)
end do



- There is an instance of S_2 that precedes an instance of S_1 in execution and S_2 consumes data that S_1 produces.
- S_2 is the source of the dependence; S_1 is the sink of the dependence.
- The dependence is loop-carried.
- The dependence distance is 1.

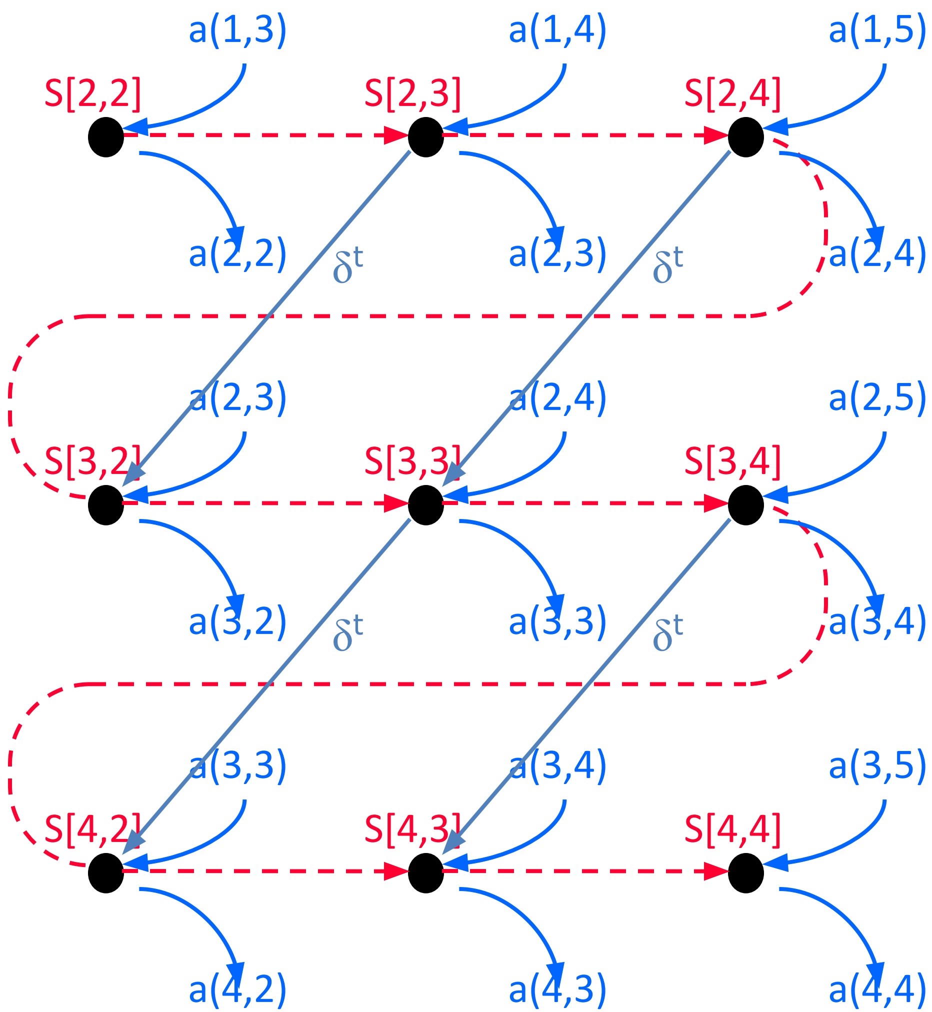
$$S_2 \, \delta^a_< \, S_1 \qquad \text{or} \qquad S_2 \, \delta^a_1 \, S_1$$

- Are you sure you know why it is $S_2 \, \boldsymbol{\delta}^a_< \, S_1$ even though S_1 appears before S_2 in the code?

Example 4

```
do i = 2, 4
  do j = 2, 4
S:    a(i,j) = a(i-1,j+1)
  end do
end do
```

- An instance of S precedes another instance of S and S produces data that S consumes.
- S is both source and sink.
- The dependence is loop-carried.
- The dependence distance is (1,-1).



$S\,\delta^t_{(<,>)}\,S$ or $S\,\delta^t_{(1,-1)}\,S$

Problem Formulation

- Consider the following perfect nest of depth d:

$$
\begin{array}{l}
\text{do } I_1 = L_1, U_1 \\
\quad \text{do } I_2 = L_2, U_2 \\
\qquad \vdots \\
\qquad \text{do } I_d = L_d, U_d \\
\qquad\quad a(f_1(\vec{I}), f_2(\vec{I}), \ldots, f_m(\vec{I})) = \ldots \\
\qquad\quad \ldots = a(g_1(\vec{I}), g_2(\vec{I}), \ldots, g_m(\vec{I})) \\
\qquad \text{enddo} \\
\qquad \vdots \\
\quad \text{enddo} \\
\text{enddo}
\end{array}
$$

$$\overbrace{a(\;\;, \underbrace{f_k(\vec{I})}_{}, \ldots, \;\;)}^{\text{array reference}}$$

subscript position

subscript function
or
subscript expression

$\vec{I} = (I_1, I_2, \ldots, I_d)$

$\vec{L} = (L_1, L_2, \ldots, L_d)$

$\vec{U} = (U_1, U_2, \ldots, U_d)$

$\vec{L} \le \vec{U}$

linear functions

$b_0 + b_1 I_1 + b_2 I_2 + \ldots + b_d I_d$

Problem Formulation

- Dependence will exist if there exists two iteration vectors $\vec{k}$ and $\vec{j}$ such that $\vec{L} \le \vec{k} \le \vec{j} \le \vec{U}$ and:

$$
\begin{array}{ll}
 & f_1(\vec{k}) = g_1(\vec{j}) \\
\text{and} & f_2(\vec{k}) = g_2(\vec{j}) \\
\text{and} & \vdots \\
\text{and} & f_m(\vec{k}) = g_m(\vec{j})
\end{array}
$$

- That is:

$$
\begin{array}{ll}
 & f_1(\vec{k}) - g_1(\vec{j}) = 0 \\
\text{and} & f_2(\vec{k}) - g_2(\vec{j}) = 0 \\
\text{and} & \vdots \\
\text{and} & f_m(\vec{k}) - g_m(\vec{j}) = 0
\end{array}
$$

Problem Formulation - Example

```
     do i = 2, 4
S1:    a(i) = b(i) + c(i)
S2:    d(i) = a(i-1)
     end do
```

- Does there exist two iteration vectors i_1 and i_2, such that
  $2 \leq i_1 \leq i_2 \leq 4$ and such that:

  $$i_1 = i_2 - 1?$$
- Answer: yes; $i_1=2$ & $i_2=3$ and $i_1=3$ & $i_2=4$.
- Hence, there is dependence!
- The dependence distance vector is $i_2-i_1 = 1$.
- The dependence direction vector is sign(1) = <.

Problem Formulation - Example

```
      do i = 2, 4
S1:     a(i) = b(i) + c(i)
S2:     d(i) = a(i+1)
      end do
```

- Does there exist two iteration vectors i_1 and i_2, such that
  $2 \leq i_1 \leq i_2 \leq 4$ and such that:

  $i_1 = i_2 + 1$?
- Answer: yes; i_1=3 & i_2=2 and i_1=4 & i_2 =3. (But, but!).
- Hence, there is dependence!
- The dependence distance vector is i_2-i_1 = -1.
- The dependence direction vector is sign(-1) = >.
- Is this possible?

Problem Formulation - Example

```
       do i = 1, 10
S1:      a(2*i) = b(i) + c(i)
S2:      d(i) = a(2*i+1)
       end do
```

- Does there exist two iteration vectors i_1 and i_2, such that $1 \leq i_1 \leq i_2 \leq 10$ and such that:

  $$2*i_1 = 2*i_2 + 1?$$

- Answer: no; $2*i_1$ is even & $2*i_2+1$ is odd.
- Hence, there is no dependence!

Problem Formulation

- Dependence testing is equivalent to an integer linear programming (ILP) problem of 2d variables & m+d constraint!
- An algorithm that determines if there exits two iteration vectors $\vec{k}$ and $\vec{j}$ that satisfies these constraints is called a dependence tester.
- The dependence distance vector is given by $\vec{j} - \vec{k}$.
- The dependence direction vector is give by sign($\vec{j} - \vec{k}$).
- Dependence testing is NP-complete!
- A dependence test that reports dependence only when there is dependence is said to be exact. Otherwise it is in-exact.
- A dependence test must be conservative; if the existence of dependence cannot be ascertained, dependence must be assumed.

Dependence Testers

- Lamport’s Test.
- GCD Test.
- Banerjee’s Inequalities.
- Generalized GCD Test.
- Power Test.
- I-Test.
- Omega Test.
- Delta Test.
- Stanford Test.
- etc…

Lamport's Test

- Lamport's Test is used when there is a single index variable in the subscript expressions, and when the coefficients of the index variable in both expressions are the same.

$$A(\cdots, b * i + c_1, \cdots) = \cdots$$
$$\cdots = A(\cdots, b * i + c_2, \cdots)$$

- The dependence problem: does there exist i_1 and i_2, such that $L_i \le i_1 \le i_2 \le U_i$ and such that

$$b*i_1 + c_1 = b*i_2 + c_2? \quad \text{or} \quad i_2 - i_1 = \frac{c_1 - c_2}{b}?$$

- There is integer solution if and only if $\frac{c_1 - c_2}{b}$ is integer.
- The dependence distance is $d = \frac{c_1 - c_2}{b}$ if $L_i \le |d| \le U_i$.
- $d > 0 \Rightarrow$ true dependence.
  $d = 0 \Rightarrow$ loop independent dependence.
  $d < 0 \Rightarrow$ anti dependence.

Lamport's Test - Example

```
do i = 1, n
  do j = 1, n
S:    a(i,j) = a(i-1,j+1)
  end do
end do
```

$i_1 = i_2 - 1$?

$b = 1;\ c_1 = 0;\ c_2 = -1$

$$\frac{c_1 - c_2}{b} = 1$$

There is dependence.

Distance (i) is 1.

$j_1 = j_2 + 1$?

$b = 1;\ c_1 = 0;\ c_2 = 1$

$$\frac{c_1 - c_2}{b} = -1$$

There is dependence.

Distance (j) is -1.

$S\,\delta^{\dagger}_{(1,-1)}\,S$ or $S\,\delta^{\dagger}_{(<,>)}\,S$

Lamport's Test - Example

```
    do i = 1, n
      do j = 1, n
S:      a(i,2*j) = a(i-1,2*j+1)
      end do
    end do
```

$i_1 = i_2 -1$?

$b = 1$; $c_1 = 0$; $c_2 = -1$

$$\frac{c_1 - c_2}{b} = 1$$

There is dependence.

Distance (i) is 1.

$2*j_1 = 2*j_2 + 1$?

$b = 2$; $c_1 = 0$; $c_2 = 1$

$$\frac{c_1 - c_2}{b} = -\frac{1}{2}$$

There is no dependence.

?

There is no dependence!

GCD Test

- Given the following equation:

$$\sum_{i=1}^{n} a_i x_i = c \qquad\qquad a_i\text{'s and } c \text{ are integers}$$

  an integer solution exists if and only if:

$$\gcd(a_1, a_2, \ldots, a_n) \quad \text{divides} \quad c$$

- Problems:
  - ignores loop bounds.
  - gives no information on distance or direction of dependence.
  - often gcd(……) is 1 which always divides c, resulting in false dependences.

GCD Test - Example

```
      do i = 1, 10
S1:     a(2*i) = b(i) + c(i)
S2:     d(i) = a(2*i-1)
      end do
```

- Does there exist two iteration vectors i_1 and i_2, such that $1 \leq i_1 \leq i_2 \leq 10$ and such that:

  $$2*i_1 = 2*i_2 - 1?$$

  or

  $$2*i_2 - 2*i_1 = 1?$$

- There will be an integer solution if and only if gcd(2,-2) divides 1.
- This is not the case, and hence, there is no dependence!

GCD Test Example

```
      do i = 1, 10
S1:     a(i) = b(i) + c(i)
S2:     d(i) = a(i-100)
      end do
```

- Does there exist two iteration vectors i_1 and i_2, such that $1 \leq i_1 \leq i_2 \leq 10$ and such that:

  $i_1 = i_2 - 100$?

  or

  $i_2 - i_1 = 100$?
- There will be an integer solution if and only if gcd(1,-1) divides 100.
- This is the case, and hence, there is dependence! Or is there?

Dependence Testing Complications

- Unknown loop bounds.

```
      do i = 1, N
S1:     a(i) = a(i+10)
      end do
```

What is the relationship between N and 10?

- Triangular loops.

```
     do i = 1, N
       do j = 1, i-1
S:       a(i,j) = a(j,i)
       end do
     end do
```

Must impose $j < i$ as an additional constraint.

More Complications

- User variables

```
    do i = 1, 10
S1:   a(i) = a(i+k)
    end do
```

```
    do i = L, H
S1:   a(i) = a(i-1)
    end do
```

Same problem as unknown loop bounds, but occur due to some loop transformations (e.g., normalization).

$\Downarrow$

```
    do i = 1, H-L
S1:   a(i+L) = a(i+L-1)
    end do
```

More Complications: Scalars

```
     do i = 1, N
S1:    x = a(i)
S2:    b(i) = x
     end do
```

⇒

```
     do i = 1, N
S1:    x(i) = a(i)
S2:    b(i) = x(i)
     end do
```

```
     j = N-1
    do i = 1, N
S1:    a(i) = a(j)
S2:    j = j - 1
     end do
```

⇒

```
    do i = 1, N
S1:    a(i) = a(N-i)

     end do
```

```
     sum = 0
    do i = 1, N
S1:    sum = sum + a(i)
    end do
```

⇒

```
    do i = 1, N
S1:    sum(i) = a(i)
    end do
    sum +=  sum(i)   i = 1, N
```

Serious Complications

- Aliases.
  - Equivalence Statements in Fortran:

    real a(10,10), b(10)

    makes b the same as the first column of a.

  - Common blocks: Fortran's way of having shared/global variables.

```
common /shared/a,b,c
          :
          :

subroutine foo (…)
common /shared/a,b,c

common /shared/x,y,z
```

Loop Parallelization

- A dependence is said to be carried by a loop if the loop is the outmost loop whose removal eliminates the dependence. If a dependence is not carried by the loop, it is loop-independent.

```
do i = 2, n-1
   do j = 2, m-1
      a(i, j) = ...
      ...          = a(i, j)

      b(i, j) = ...
      ...     = b(i, j-1)

      c(i, j) = ...
      ...     = c(i-1, j)
   end do
end do
```

Loop Parallelization

A dependence is said to be carried by a loop if the loop is the outmost loop whose removal eliminates the dependence. If a dependence is not carried by the loop, it is loop-independent.

$\delta^{\dagger}_{=,=}$

```
do i = 2, n-1
   do j = 2, m-1
      a(i, j) = ...
       ...            = a(i, j)

      b(i, j) = ...
       ...      = b(i, j-1)

      c(i, j) = ...
       ...      = c(i-1, j)
   end do
end do
```

Loop Parallelization

A dependence is said to be carried by a loop if the loop is the outmost loop whose removal eliminates the dependence. If a dependence is not carried by the loop, it is loop-independent.

```
do i = 2, n-1
   do j = 2, m-1
      a(i, j) = ...
       ...          = a(i, j)

      b(i, j) = ...
       ...     = b(i, j-1)

      c(i, j) = ...
       ...     = c(i-1, j)
   end do
end do
```

$\delta^{t}_{=,<}$

Loop Parallelization

A dependence is said to be carried by a loop if the loop is the outmost loop whose removal eliminates the dependence. If a dependence is not carried by the loop, it is loop-independent.

```
do i = 2, n-1
   do j = 2, m-1
      a(i, j) = ...
      ...          = a(i, j)

      b(i, j) = ...
      ...     = b(i, j-1)

      c(i, j) = ...
      ...     = c(i-1, j)
   end do
end do
```

$\delta^{t}_{<,=}$

Loop Parallelization

A dependence is said to be carried by a loop if the loop is the outmost loop whose removal eliminates the dependence. If a dependence is not carried by the loop, it is loop-independent.

```
do i = 2, n-1
   do j = 2, m-1
      a(i, j) = ...
       ...              = a(i, j)

      b(i, j) = ...
       ...     = b(i, j-1)

      c(i, j) = ...
       ...     = c(i-1, j)
   end do
end do
```

$\delta^{\dagger}_{=,=}$ (a), $\delta^{\dagger}_{=,<}$ (b), $\delta^{\dagger}_{<,=}$ (c)

- Outermost loop with a non "=" direction carries dependence!

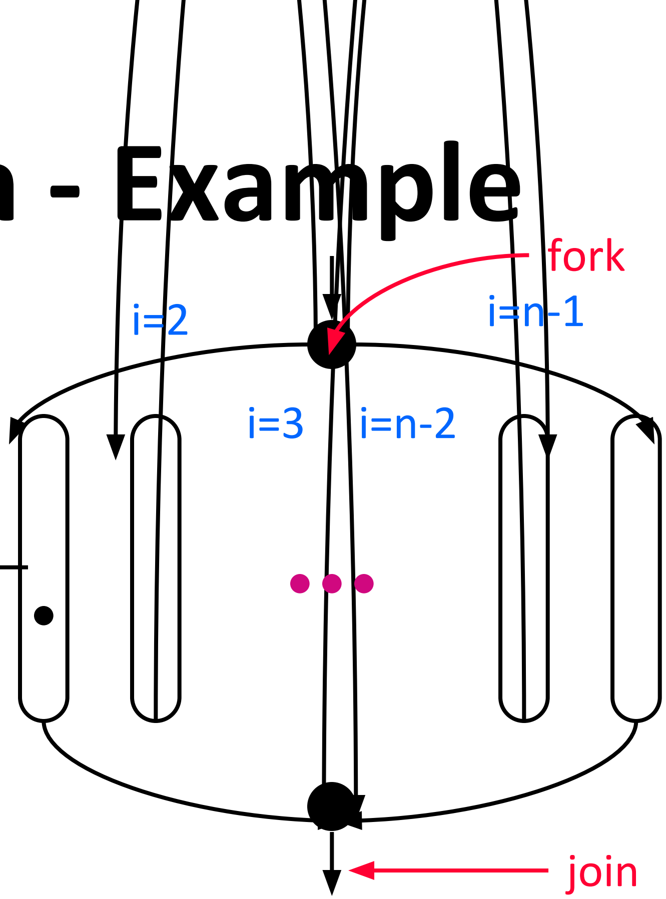
Loop Parallelization

The iterations of a loop may be executed in parallel with one another if and only if no dependences are carried by the loop!

Loop Parallelization - Example

$\delta^{t}_{=,<}$

```
do i = 2, n-1
    do j = 2, m-1
       b(i, j)      = ...
         ...    = b(i, j-1)
    end do
end do
```

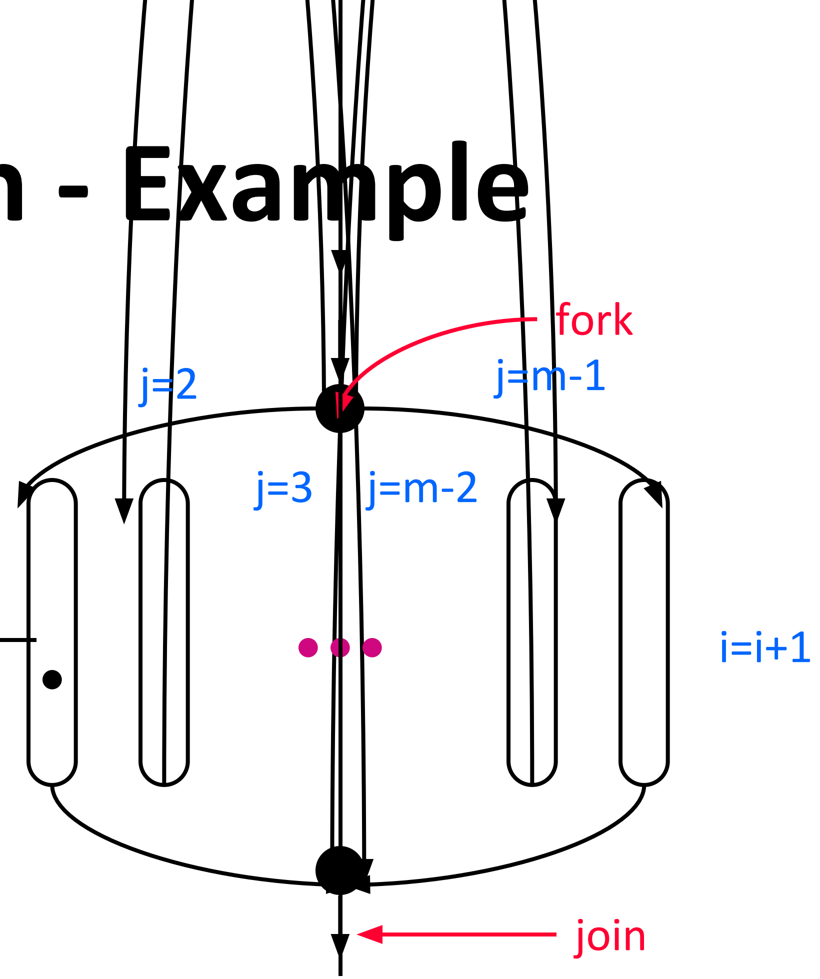


- Iterations of loop j must be executed sequentially, but the iterations of loop i may be executed in parallel.
- Outer loop parallelism.

Loop Parallelization - Example

$\delta^{\dagger}_{<,=}$

```
do i = 2, n-1
   do j = 2, m-1
      b(i, j)      = ...
      ...    = b(i-1, j)
   end do
end do
```

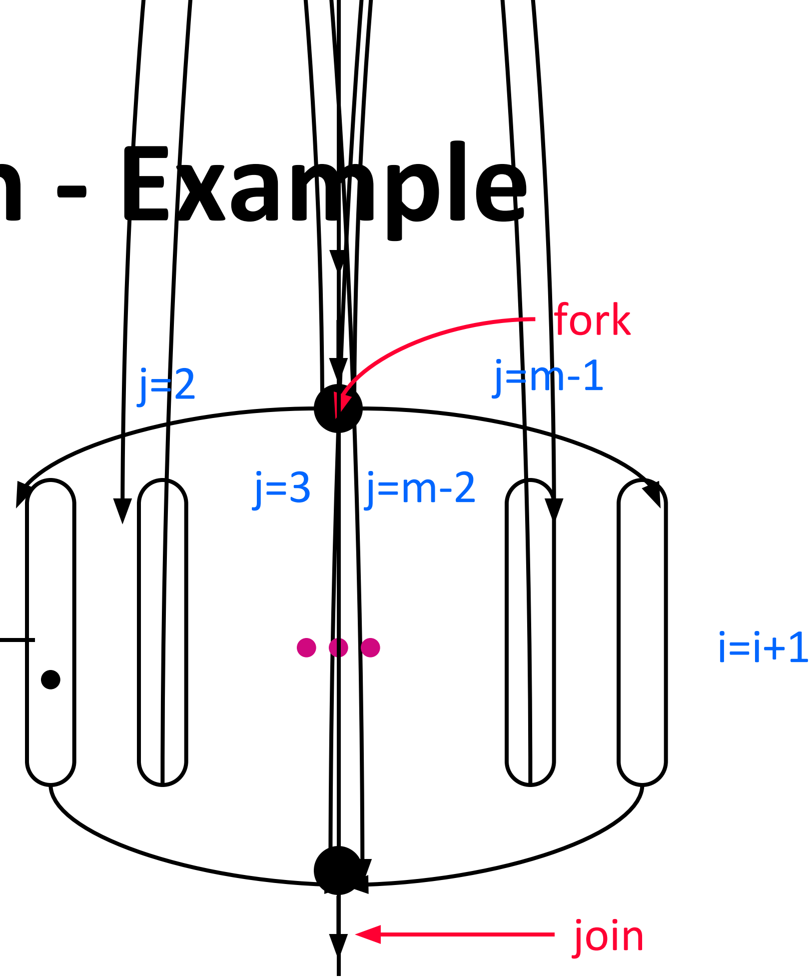


- Iterations of loop i must be executed sequentially, but the iterations of loop j may be executed in parallel.
- Inner loop parallelism.

Loop Parallelization - Example

$\delta^{t}_{<,<}$

```
do i = 2, n-1
   do j = 2, m-1
      b(i, j)      = ...
         ...  = b(i-1, j-1)
   end do
end do
```

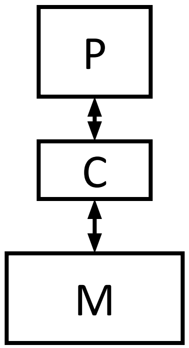


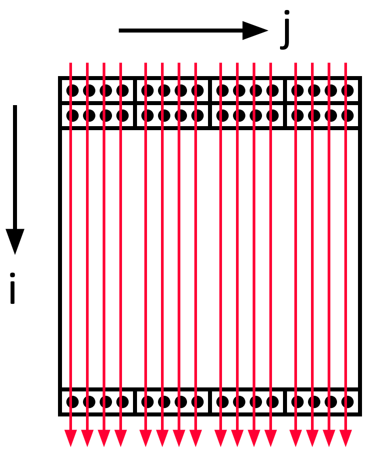
- Iterations of loop i must be executed sequentially, but the iterations of loop j may be executed in parallel. Why?
- Inner loop parallelism.

Loop Interchange

Loop interchange changes the order of the loops to improve the spatial locality of a program.

```
do j = 1, n
  do i = 1, n
       ... a(i,j) ...
  end do
end do
```



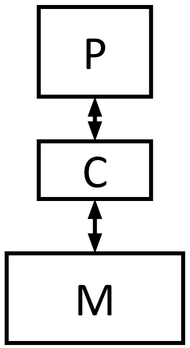


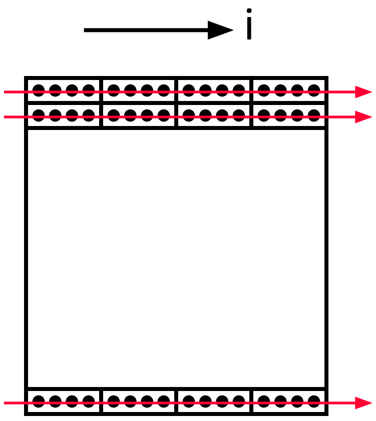
Loop Interchange

Loop interchange changes the order of the loops to improve the spatial locality of a program.

```
do j = 1, n
  do i = 1, n
     ... a(i,j) ...
  end do
end do
```

```
do i = 1, n
  do j = 1, n
     ... a(i,j) ...
  end do
end do
```





Loop Interchange

- Loop interchange can improve the granularity of parallelism!

```
do i = 1, n
  do j = 1, n
     a(i,j) = b(i,j)
     c(i,j) = a(i-1,j)
  end do
end do
```

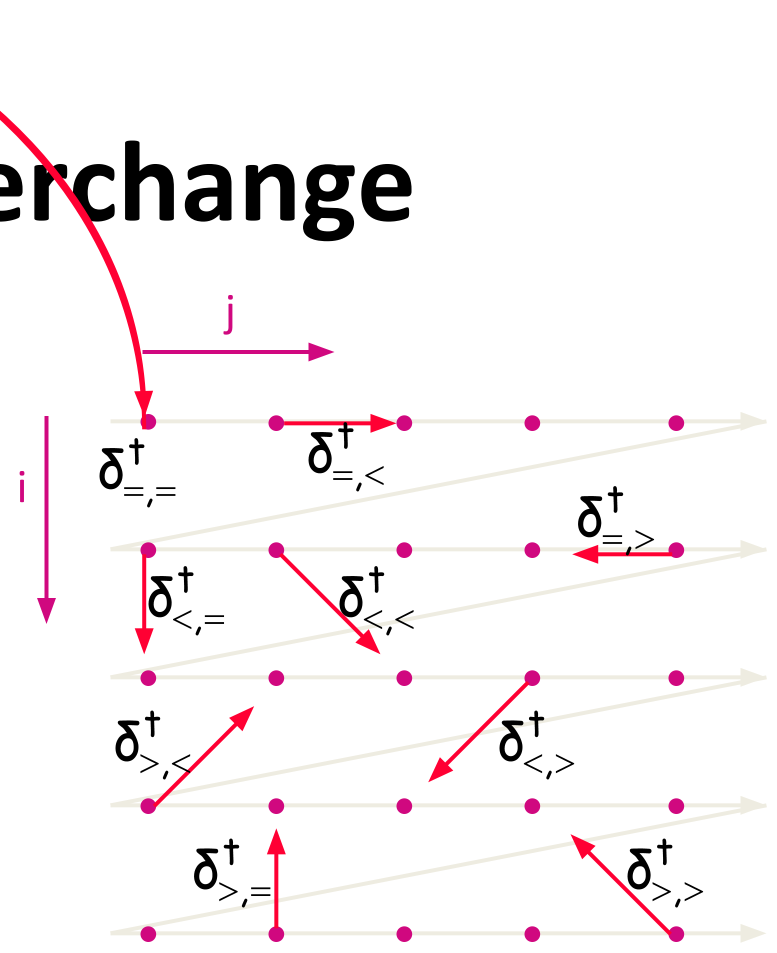
```
do j = 1, n
  do i = 1, n
     a(i,j) = b(i,j)
     c(i,j) = a(i-1,j)
  end do
end do
```

$\delta^{\dagger}_{<,=}$

$\delta^{\dagger}_{=,<}$

Loop Interchange

do i = 1,n
  do j = 1,n
    … a(i,j) …
  end do
end do

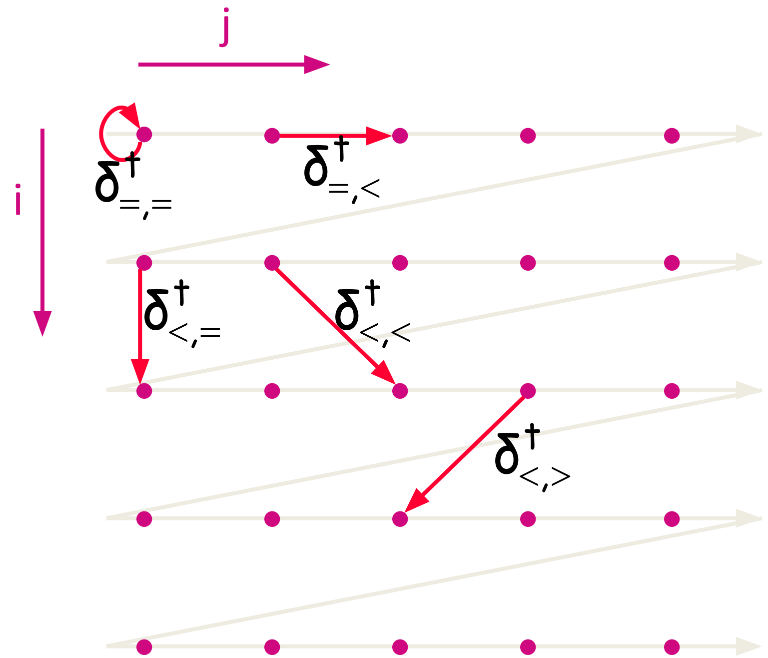


do j = 1,n
  do i = 1,n
    … a(i,j) …
  end do
end do

- When is loop interchange legal?

Loop Interchange

```
do i = 1,n
  do j = 1,n
    ... a(i,j) ...
  end do
end do
```

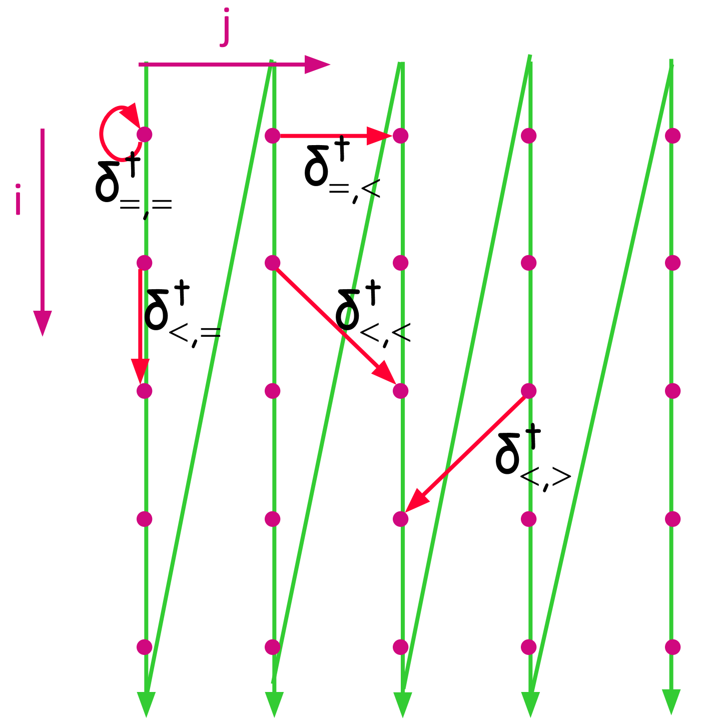



```
do j = 1,n
  do i = 1,n
    ... a(i,j) ...
  end do
end do
```

- When is loop interchange legal?

Loop Interchange

```
do i = 1,n
  do j = 1,n
    ... a(i,j) ...
  end do
end do
```

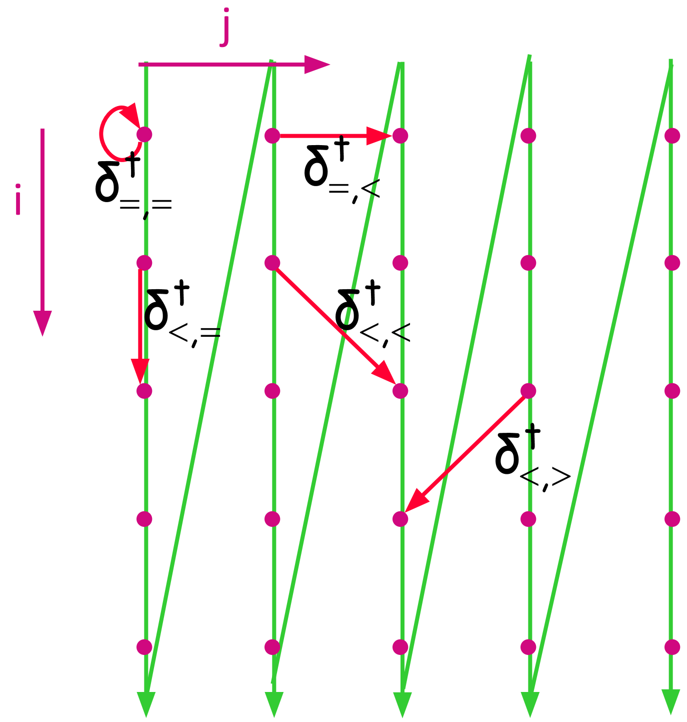



```
do j = 1,n
  do i = 1,n
    ... a(i,j) ...
  end do
end do
```

- When is loop interchange legal?

Loop Interchange

```
do i = 1,n
  do j = 1,n
     ... a(i,j) ...
  end do
end do
```



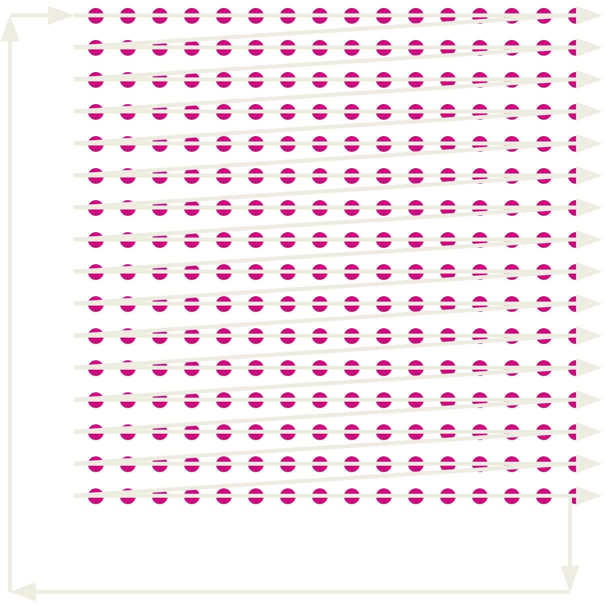

```
do j = 1,n
  do i = 1,n
     ... a(i,j) ...
  end do
end do
```

- When is loop interchange legal? when the “interchanged” dependences remain lexiographically positive!

Loop Blocking (Loop Tiling)

Exploits temporal locality in a loop nest.

```
do t = 1,T
  do i = 1,n
    do j = 1,n
       … a(i,j) …
    end do
  end do
end do
```

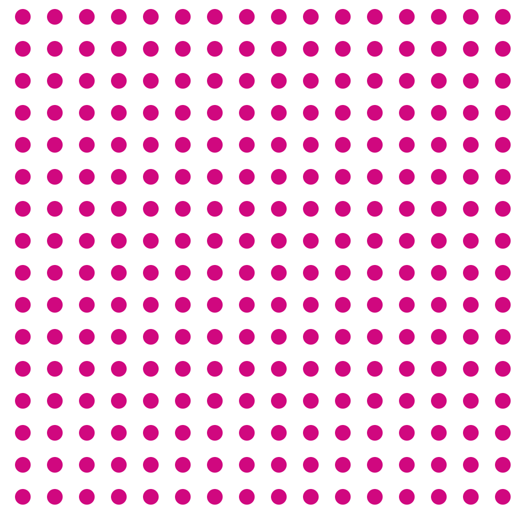


Loop Blocking (Loop Tiling)

Exploits temporal locality in a loop nest.

control loops

```
do ic = 1, n, B
  do jc = 1, n , B
    do t = 1,T
      do i = 1,B
        do j = 1,B
          … a(ic+i-1,jc+j-1) …
        end do
      end do
    end do
  end do
end do
```



B: Block size

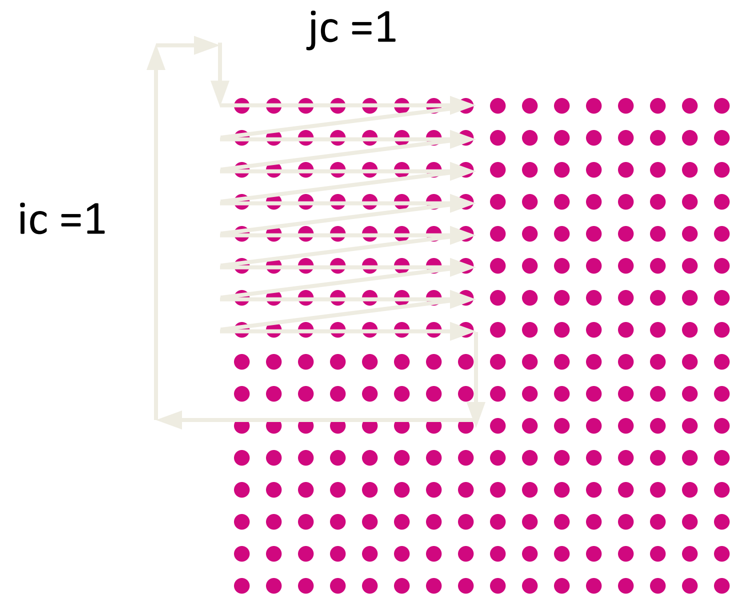
Loop Blocking (Loop Tiling)

Exploits temporal locality in a loop nest.

```
do ic = 1, n, B
  do jc = 1, n , B
    do t = 1,T
      do i = 1,B
        do j = 1,B
          … a(ic+i-1,jc+j-1) …
        end do
      end do
    end do
  end do
end do
```

control loops

B: Block size



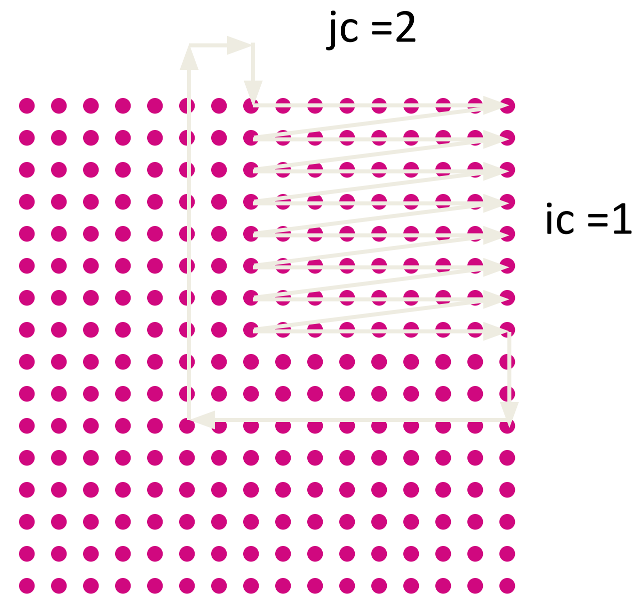
Loop Blocking (Loop Tiling)

Exploits temporal locality in a loop nest.

```
do ic = 1, n, B
  do jc = 1, n , B
    do t = 1,T
      do i = 1,B
        do j = 1,B
          … a(ic+i-1,jc+j-1) …
        end do
      end do
    end do
  end do
end do
```

control loops

B: Block size



Loop Blocking (Loop Tiling)

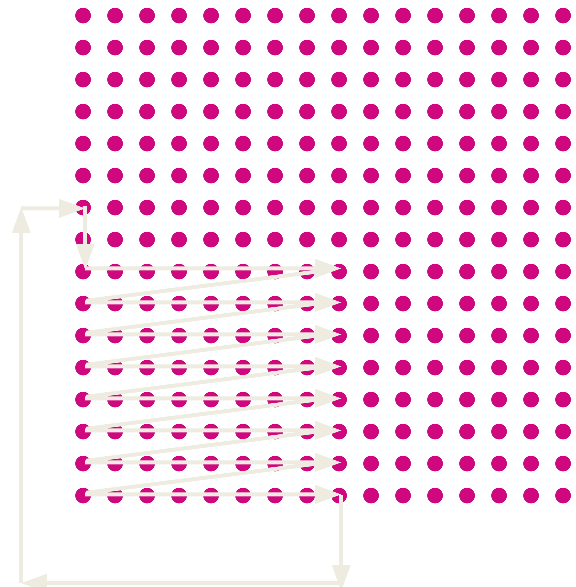
Exploits temporal locality in a loop nest.

control loops

```
do ic = 1, n, B
  do jc = 1, n , B
    do t = 1,T
      do i = 1,B
        do j = 1,B
          … a(ic+i-1,jc+j-1) …
        end do
      end do
    end do
  end do
end do
```

B: Block size

ic =2



jc =1

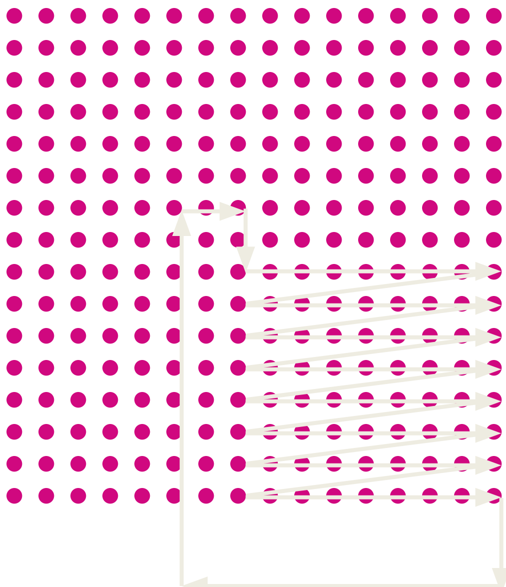
Loop Blocking (Loop Tiling)

Exploits temporal locality in a loop nest.

```
do ic = 1, n, B
  do jc = 1, n , B
    do t = 1,T
      do i = 1,B
        do j = 1,B
          … a(ic+i-1,jc+j-1) …
        end do
      end do
    end do
  end do
end do
```

control loops

B: Block size



Loop Blocking (Tiling)

```
do t = 1,T
  do i = 1,n
    do j = 1,n
       … a(i,j) …
    end do
  end do
end do
```

```
do t = 1,T
 do ic = 1, n, B
  do i = 1,B
    do jc = 1, n, B
     do j = 1,B
        … a(ic+i-1,jc+j-1) …
     end do
  end do
 end do
end do
```

```
do ic = 1, n, B
  do jc = 1, n , B
     do t = 1,T
        do i = 1,B
           do j = 1,B
              … a(ic+i-1,jc+j-1) …
           end do
        end do
     end do
  end do
end do
```

- When is loop blocking legal?

CSC D70: Compiler Optimization Parallelization

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University of Toronto

Winter 2019

The content of this lecture is adapted from the lectures of Todd Mowry and Tarek Abdelrahman